

Confidential QUBO solver

Mariano Caruso^{1,2,3,4} , Daniel Escanez-Exposito⁵ , Pino Caballero-Gil⁵ ,
and Carlos Kuchkovsky⁴ 

¹ UGR, Granada, Spain.

`mcaruso@ugr.es`

² FIDESOL, Granada, Spain.

³ UNIR, Logroño, Spain.

⁴ QCentroid, Bilbao, Spain.

⁵ CryptULL–ULL, Tenerife, Spain.

Abstract. Quadratic Unconstrained Binary Optimization (QUBO) is widespread and solvable via classical or quantum computing. However, outsourcing these computations online exposes sensitive data to potential breaches. We introduce a novel encryption scheme that seamlessly integrates with any solver or hardware platform, ensuring data security without compromising performance. A robust `python` implementation delivers promising results, marking a significant step forward in secure optimization for both classical and quantum environments.

Keywords: QUBO problems, secure optimization, cryptographic solution, transfer principle, quantum computing, hardware agnostic.

1 Introduction

Quadratic Unconstrained Binary Optimization (QUBO) is crucial in combinatorial optimization across disciplines like cryptography, economics, physics, and machine learning [1–8]. It models problems such as max-cut, graph coloring, and clustering [9–11, 13], with applications in quantum computing via Ising models and quantum annealing [14–16]. As NP-hard [17], QUBO demands efficient and secure solving methods.

Online solvers risk exposing sensitive data. This work proposes an encryption scheme enabling secure external solving without revealing original data. Unlike standard homomorphic encryption [18–20], our approach meets the specific needs of QUBO.

We define the set of natural indices $I_n := \{1, \dots, n\}$ and the binary space $\{0, 1\}^n$. A function $f(\mathbf{x}) = \sum_{(i,j) \in I_n^2} \mathbf{Q}_{ij} x_i x_j$ is defined using a matrix $\mathbf{Q} \in \mathbb{R}^{n \times n}$. Since adding a constant does not change the optimal solution, the typical QUBO problem is to find

$$\underset{\mathbf{x} \in \{0,1\}^n}{\operatorname{argmin}} f(\mathbf{x}),$$

where the objective function f can be expressed compactly as $f(\mathbf{x}) = \mathbf{x} \cdot \mathbf{Q} \mathbf{x}$. When solving such problems with cloud solvers, transmitting the matrix \mathbf{Q} in clear text risks exposing sensitive information.

2 Transfer principle

The encrypted objects are denoted with a prime, so the encrypted objective function is $f'(\mathbf{x}') = \mathbf{x}' \cdot \mathbf{Q}' \mathbf{x}'$. Define a mapping T from the original optimization problem to an encrypted version $f'(\mathbf{x}') := f(T(\mathbf{x}'))$. This allows a user to send an encrypted QUBO problem to an external solver without exposing sensitive data (the matrix \mathbf{Q}). The transformation T defines a new optimization problem in such a way that its resolution in the encrypted domain can be carried out without revealing information about the original optimization problem and without requiring key exchange. In this context, we could mention some similarities with homomorphic encryption. Note that T will not be a homomorphism in all cases, as it will not preserve the operations of addition \oplus or product \odot in $\{0, 1\}^n$. With a transformation T satisfying $\mathbf{x} = T(\mathbf{x}')$, we establish a relation between the original matrix \mathbf{Q} and its encrypted counterpart \mathbf{Q}' . Matching the quadratic forms in the expressions for f and f' yields a closed-form relation through an encryption function: $\mathbf{Q}' = \mathcal{E}(\mathbf{k}_T, \mathbf{Q})$ or $\mathbf{Q}' = \mathcal{E}_{\mathbf{k}_T}(\mathbf{Q})$ [21]. The matrix of the original QUBO problem is encrypted and sent to the `QUBO – solver`. The solution returned, \mathbf{x}'_* , is then decrypted locally, and the solution to the original problem is obtained via $\mathbf{x}_* = T(\mathbf{x}'_*)$. In order to construct T there are two ingredients: diffusion via permutations and confusion via substitutions, and the encrypted matrix \mathbf{Q}' is obtained from $\mathcal{E}_{\mathbf{P}}$ and $\mathcal{E}_{\mathbf{k}}$, respectively.

3 Conclusions

This work addresses the security challenges in the online resolution of unconstrained binary quadratic optimization problems. By employing cryptographic methods, we establish a framework for securely solving these problems using both classical and quantum computing. This approach protects the problem’s sensitive information during resolution and encourages the adoption of advanced techniques in online environments, promoting further development in quantum computing and data security.

A complete implementation of these proposals has been developed in `python`. The proposed cryptographic solution is independent of the optimization method, so it does not depend on how the QUBO problem is solved. It supports exact methods, classical or quantum annealing, as well as heuristic approaches such as simulated annealing and genetic algorithms, making it hardware agnostic.

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